Section 1.3: Quantifiers, Predicates, and Validity

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Abstract

We now throw variables into the mix, and investigate wffs which describe properties of the domains of those variables in given "interpretations." We still test their truth values, either for the specific domain in question, or even in all domains (validity).

1 Predicates and quantifiers

• quantifier: tells how many objects in a given domain have a certain property.

Examples:

- universal quantifier \forall "for all", "for every", "for any"
- existential quantifier $-\exists$ "there exists", "for at least one", "for some"

Have you encountered these quantifiers before, in other courses?

• **predicate**: a property of a variable (e.g. "x is prime"), generally containing one or more variables (and perhaps some constants).

We combine the quantifiers and predicates to create expressions (wffs) such as

$$(\forall x)P(x)$$

which we then must *interpret*. For example, this might be said in the context of the integers, with P(x) standing for "x is prime". (So this wff would be false in this context.)

There is nothing special about the variable x, so this wff is the same as $(\forall y)P(y)$, $(\forall z)P(z)$, etc. We say that x is a dummy variable.

Predicates may have any number of variables in them: the example above is a *unary* predicate, with only a single variable.

• Truth value hence now depends on the **Interpretation** of an expression:

- domain of interpretation non-empty set to which the predicate expression is applied
- assignment of a property of the objects to each predicate in the expression
- assignment of particular objects to each constant symbol in the expression

We start with something abstract, and replace it with concrete instances in a given context.

Example: Ex.
$$2(e, f)$$
, p. $42/43$
 $(4\times)(4 y)(Q(x,y) \times P(x,y))$
 $2e. (4x)[4y)(x < y \times y < x)$
 $false : \times could$
 $equal y$
 $2f$, $(4x)[x < 0 \rightarrow (3y)(y > 0 \land x + y = 0)]$
 $false : Could$
 $false : Could$

Example: Ex. 4/5(b, d), p. 42/44

45.
$$(4x)(4y)[P(x,y) \rightarrow P(y,x)]$$

Domain: Z $P(x,y): "x = y"$

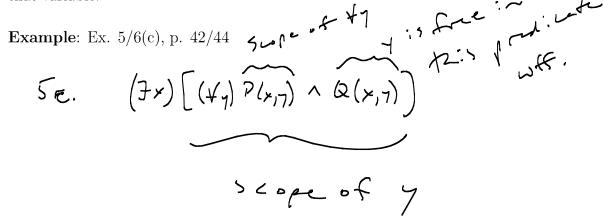
Truc in $TZ:S$ intropretation

Same domain: $P(x,y): "x = y"$

False in $TZ:S$ one.

In example 4/5b above, the predicate $P(x,y) \longleftrightarrow P(y,x)$ is an example of a binary predicate.

- scope of a quantifier: the part of an expression to which the quantifier applies; indicated with parentheses or brackets (but these may be neglected if the scope is clear).
- free variable: a variable in a predicate wff outside the scope of a quantifier involving that variable.



Translation of English statements into predicate wffs 2

As noted before, this can be a very tricky business, but an important one. As is often the case, the process of translation does not result in a unique expression: there may be several different ways to do the same job.

Our author encourages us to remember that

- typically \exists and \land go together, whereas
- typically \forall and \rightarrow go together.

Also, a single English statement may be given by numerous wffs.

Example: Ex. 9/14(e, g), p. 43/46

Only judges admire judges "Some women admire

T(x) - x 11 a judge (7x)(4x) [11x] A(x,y) - x admins y

Example: Ex. 11/16(a-d), p. 44/47

c.
$$(\forall y) [\forall y) [\exists (x) \land F(y) \Rightarrow L(x,y)]$$

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$$(\forall y) [\exists (x) \Rightarrow (\forall y) (F(y) \Rightarrow L(x,y)]$$

$$(\forall y) [\exists (x) \Rightarrow (\forall y) (F(y) \Rightarrow L(x,y))]$$

$$(\forall y) [\exists (x) \land F(y) \Rightarrow L(x,y)]$$

$$(\exists x) [\exists (x) \land (\forall y) [F(y) \Rightarrow L(x,y)]$$

$$(\exists x) [\exists (x) \land (\forall y) [F(y) \Rightarrow L(x,y)]$$

• Negation of predicate wffs: some cases are standard, e.g.

c. (Yx) [B(x) -> (3y) [F(7) 1 L(x,4)])

- The negation of "Every x has property A." is "There is an x which doesn't have property A." $[(\forall x)A(x)]' \iff (\exists x)[A(x)]'$
- The negation of "There is an x which has property A." is "No x has property A." $[(\exists x)A(x)]'\iff (\forall x)[A(x)]'$

In general, English makes negation kind of tricky. Watch your step!

Example: Ex. 14/19(c,d), p. 44/48

3 Validity

The truth value of a predicate wff depends on the interpretation, but there are some for which the wff is true independent of the interpretation. These are called **valid** predicate wffs (the analogue of tautology for propositional wffs).

Whereas we can check the "validity" of a propositional wff (just check the truth table to see if it's a tautology), there is no general check for the validity of a predicate wff, since it depends on context. In spite of that, there are some valid predicate wffs (context free truth!), as demonstrated in the text:

$$(\forall x)P(x) \to (\exists x)P(x)$$
$$(\forall x)P(x) \to P(a)$$
$$P(x) \to (Q(x) \to P(x))$$

Example: Ex. 18/24(d,e), p. 45/49

e.
$$(4x) \lceil A(x) \rightarrow (34) \rceil \rightarrow (4x) A(x) \rightarrow (4x) 360 \rceil$$

More on 9g: Re first one we tried didn't work, a neiter does this.

(3x) [(47) [U(x) \(A (\times x)) -> L(y)']]

No - anything in the domain that's not a woman makes this true. Same with this one,

(3x) [(4y) [U(x) \(L(y) -> A(x,y)']]

and some with owr in itial version'.

"Some women admire no lawyer." (Jx)(Hy)[W(x) \(\lambda\) \tagentally) \tagentally

All of these fail to capture the fact that we are asserting the existence of a woman who admires no lawyers.

(fx)[W&)~....]

I got back on track at the board only for a moment when I looked at one of these of said "A man makes it true!" But someone throw me off my game by asserting that "Yes, but we've said there's a woman, too"

Ratis correct, but the fact of the matter is that all thru of these with are true even it no woman admires only lawyers!

Because a man will make it true, regardless.

And Rat is a bad thing....



10.f If the is a Corrette that 11 slower than a Ferri, the all Correttes are slower than all Ferris.

$$(J_{\times})[C(x) \wedge (J_{\gamma})(F(\gamma) \wedge S(x,\gamma)] \longrightarrow$$

$$(J_{\times})(J_{\gamma})[C(x) \wedge F(\gamma) \rightarrow S(x,\gamma)]$$

20/26 c