MARIE (chapter 4) sample problems
Provide the full RTN for a new MARIE instruction, Inc X. This instruction fetches datum X from memory, increments it and stores it back to memory.

```
Fetch: \(\quad\) MAR \(\leftarrow P C\)
    MBR \(\leftarrow \mathrm{M}[\mathrm{MAR}] \quad / /\) divided into two steps
    \(\mathrm{IR} \leftarrow \mathrm{MBR}\)
    \(\mathrm{PC} \leftarrow \mathrm{PC}+1\)
Decode: Decode IR[15..12] // reversed the order of these two steps
    MAR \(\leftarrow \operatorname{IR}[11 . .0]\)
Get op: \(\quad \mathrm{MBR} \leftarrow \mathrm{M}[\mathrm{MAR}]\)
Execute: \(\quad \mathrm{AC} \leftarrow \mathrm{MBR}+1 \quad / /\) increment X
    \(\mathrm{MBR} \leftarrow \mathrm{AC} \quad / /\) move \(\mathrm{X}+1\) back to MBR
    \(\mathrm{M}[\mathrm{MAR}] \leftarrow \mathrm{MBR} \quad / /\) store back to memory location X
```

Provide the full RTN for a new MARIE instruction, JZ X. This instruction sets the AC to see if it is currently 0 and if so, branches to memory location X .

```
Fetch: \(\quad\) MAR \(\leftarrow\) PC
    MBR \(\leftarrow \mathrm{M}[\mathrm{MAR}]\)
    \(\mathrm{IR} \leftarrow \mathrm{MBR}\)
    \(\mathrm{PC} \leftarrow \mathrm{PC}+1\)
Decode: Decode IR[15..12]
    MAR \(\leftarrow \operatorname{IR}[11 . .0] \quad / /\) this step is not needed
Get op: not needed
Execute: If \(\mathrm{AC}=0\) then \(\mathrm{PC} \leftarrow \operatorname{IR}[11 . .0] / /\) if we do the above step of
    \(/ /\) MAR \(\leftarrow \operatorname{IR}[11 . .0]\), we could do \(\mathrm{PC} \leftarrow\) MAR here
```

We decide to expand MARIE to have 4 registers R0, R1, R2, R3, instead of using the AC. Our instructions now must include the register being referenced as in Load R1, X or Add R3, Y. What impact will this change have on all of MARIE? Be as complete as possible (note: there are two ways you could go with this answer).

Both answers are based on the fact that to denote the register, we need 2 bits.
Answer 1: We add 2 bits to our instruction so that all instructions are now 18 bits long. This will require that the word size expand from 16 bits to 18 bits so that memory is now 4 Kx 18 instead of 4 Kx 16 . It will require an 18 bit data bus and 18 bit sized registers for the MBR, IR and the 4 data registers. It will also require expanding the ALU to operate on 18 bits.

Answer 2: We remove 2 bits from the operand specification meaning that instead of 12 bits for a memory address, we only use 10 . This causes our address space to shrink from

4 K to 1 K . We would also reduce the address bus from 12 bits to 10 and the MAR and PC from 12 bits to 10 .

Provide the MARIE code for the corresponding pseudocode below (written in a Java-like format except for input and output).

```
input(x, y, z);
if(x>y&&y>z)
    output(x);
else if(y!=z) output(y);
else output(0);
Input
Store X
Input
Store Y
Input
Store Z
Load X
Subt Y
Skipcond 10
Jump elseif
Load Y
Subt Z
Skipcond 10
Jump elseif
Load X
Output
Jump done
Subt Z
Skipcond 01 // true if y==z
Jump outy // we reach here if the above condition is false, thus we want to output y
Load #0
Output
Jump done
    Output
...
```

elseif: Load Y
outy: Load Y
done:

Write MARIE code to compute $\mathrm{z}=\mathrm{x} / \mathrm{y}$. Assume its ok to destroy the contents of x in doing so.
Clear
Store z
top: $\quad \operatorname{Load} \mathrm{x}$
Subt y
Skipcond 10 // having subtracted y from x , is x still positive?
Jump done
Store $\mathrm{x} \quad / /$ if so, we continue
Load z
Add \#1
Store z
Jump top
done: ... // note: z is the quotient and x is the remainder

Write the following pseudocode into MARIE code:
Input $x$
While ( $\mathrm{x}>=\mathrm{y}$ ) $\operatorname{If}(x \% 2==0) y++;$
Input x
Output y
NOTE: assume we have a subroutine called $\bmod 2$ which computes $\mathrm{x} \% 2$ and sets the variable temp to 1 if the result is equal to 0 .

Input
Store x
top: $\operatorname{Load} x$
Subt y
Skipcond 00 // we do not have a test for $>=$, so I'm using $<$
Jump body // if $x>=y$, we go here, so we jump to the load body
Jump done // otherwise we skip to here and we want to skip the loop body
Clear
Store temp
Jns mod2
Load temp
Subt \#1 $\quad / /$ does temp $==1$ ?
Skipcond 01
Jump bottom // since temp != 1, $\mathrm{x} \% 2$ != 0 , go to the bottom of the loop body
Load y
Add \#1
Store y
bottom: Input
Store x
Jump top
done: Load y
Output

Hand-compile into hexadecimal the program in example 4.3 on page 253. NOTE: Skipcond 400 should be Skipcond 01.

| 100 | Load X | 110 C |
| :--- | :--- | :--- |
| 101 | Subt Y | 410 D |
| 102 | Skipcond 01 | 8400 |
| 103 | Jump Else | 9108 |
| 104 | Load X | 110 C |
| 105 | Add X | 310 C |
| 106 | Store X | 210 C |
| 107 | Jump Endif | 910 B |
| 108 | Load Y | 110 D |
| 109 | Subt X | 410 C |
| 10 A | Store Y | 210 D |
| $10 B$ | Halt | 7000 |

