

THE DISCRETE HAAR WAVELET TRANSFORMATION

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Joint Mathematical Meetings, 7 & 9 January 2008



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- ▶ Suppose you are given N values

$$\mathbf{x} = (x_1, x_2, \dots, x_N)$$

where N is even.

- ▶ **Your task:** Send an approximation \mathbf{s} (a list of numbers) of this data via the internet to a colleague.
- ▶ In order to reduce transfer time, the length of your approximation must be $N/2$.
- ▶ How do you suggest we do it?



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- ▶ One solution is to pair-wise average the numbers:

$$s_k = \frac{x_{2k-1} + x_{2k}}{2}, \quad k = 1, \dots, N/2$$

- ▶ For example:

$$\mathbf{x} = (6, 12, 15, 15, 14, 12, 120, 116) \rightarrow \mathbf{s} = (9, 15, 13, 118)$$



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- ▶ Suppose now you were allowed to send extra data in addition to the pair-wise averages list \mathbf{s} .
- ▶ The idea is to send a second list of data \mathbf{d} so that the original list \mathbf{x} can be recovered from \mathbf{s} and \mathbf{d} .
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- ▶ There are a couple of choices for d_k (called **directed distances**):
- ▶ We could set

$$d_k = \frac{x_{2k-1} - x_{2k}}{2}, \quad k = 1, \dots, N/2$$

- ▶ or

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- ▶ The process is invertible since

$$s_k + d_k = \frac{x_{2k-1} + x_{2k}}{2} + \frac{x_{2k} - x_{2k-1}}{2} = x_{2k}$$

and

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- ▶ So we map $\mathbf{x} = (x_1, x_2, \dots, x_N)$ to $(\mathbf{s} \mid \mathbf{d}) = (s_1, \dots, s_{N/2} \mid d_1, \dots, d_{N/2})$.
- ▶ Using our example values we have

$$(6, 12, 15, 15, 14, 12, 120, 116) \rightarrow (9, 15, 13, 118 \mid 3, 0, -1, -2)$$

- ▶ Why might people prefer the data in this form?



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- ▶ We can identify large changes in the the differences portion \mathbf{d} of the transform.
- ▶ It is easier to quantize the data in this form.
- ▶ The transform concentrates the signal's energy in fewer values.
- ▶ And the obvious answer: **less digits!!**
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- ▶ The transformation

$$\mathbf{x} = (x_1, \dots, x_N) \rightarrow (\mathbf{s} \mid \mathbf{d}) = (s_1, \dots, s_{N/2} \mid d_1, \dots, d_{N/2})$$

is called the **Discrete Haar Wavelet Transformation**.

- ▶ What does the transform look like as a matrix?



- ▶ The transformation

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- ▶ What does the transform look like as a matrix?



Consider applying the transform to an 8-vector. What is the matrix that works?

$$\begin{bmatrix} \\ \\ \\ \\ \\ \\ \\ \end{bmatrix} \cdot \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \\ x_7 \\ x_8 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} x_1 + x_2 \\ x_3 + x_4 \\ x_5 + x_6 \\ x_7 + x_8 \\ \hline x_2 - x_1 \\ x_4 - x_3 \\ x_6 - x_5 \\ x_8 - x_7 \end{bmatrix}$$



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$$\begin{bmatrix}
 \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\
 \hline
 -\frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 & 0 & 0 \\
 0 & 0 & -\frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 \\
 0 & 0 & 0 & 0 & -\frac{1}{2} & \frac{1}{2} & 0 & 0 \\
 0 & 0 & 0 & 0 & 0 & 0 & -\frac{1}{2} & \frac{1}{2}
 \end{bmatrix} \cdot \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \\ x_7 \\ x_8 \end{bmatrix} = \frac{1}{2} \begin{bmatrix} x_1 + x_2 \\ x_3 + x_4 \\ x_5 + x_6 \\ x_7 + x_8 \\ \hline x_2 - x_1 \\ x_4 - x_3 \\ x_6 - x_5 \\ x_8 - x_7 \end{bmatrix}$$

We will denote the transform matrix by W_8 .



What about W_8^{-1} ? That is, what matrix solves

$$\begin{bmatrix} \\ \\ \\ \\ \\ \\ \\ \end{bmatrix} \cdot \left(\frac{1}{2} \begin{bmatrix} x_1 + x_2 \\ x_3 + x_4 \\ x_5 + x_6 \\ x_7 + x_8 \\ \hline x_2 - x_1 \\ x_4 - x_3 \\ x_6 - x_5 \\ x_8 - x_7 \end{bmatrix} \right) = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \\ x_7 \\ x_8 \end{bmatrix}$$



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$$\begin{bmatrix} 1 & 0 & 0 & 0 & -1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & -1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & -1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & -1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \end{bmatrix} \cdot \left(\frac{1}{2} \begin{bmatrix} X_1 + X_2 \\ X_3 + X_4 \\ X_5 + X_6 \\ X_7 + X_8 \\ \hline X_2 - X_1 \\ X_4 - X_3 \\ X_6 - X_5 \\ X_8 - X_7 \end{bmatrix} \right) = \begin{bmatrix} X_1 \\ X_2 \\ X_3 \\ X_4 \\ X_5 \\ X_6 \\ X_7 \\ X_8 \end{bmatrix}$$



- ▶ Learning how to code the HWT and its inverse provides a good review of linear algebra.
- ▶ We'll use $N = 8$ as an example. Let

$$W_8 = \begin{bmatrix} \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\ \hline -\frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & -\frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & -\frac{1}{2} & \frac{1}{2} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & -\frac{1}{2} & \frac{1}{2} \end{bmatrix} = \begin{bmatrix} H \\ - \\ G \end{bmatrix}$$



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► Then

$$W_8 \mathbf{x} = \begin{bmatrix} H \\ - \\ G \end{bmatrix} \mathbf{x} = \begin{bmatrix} H\mathbf{x} \\ - \\ G\mathbf{x} \end{bmatrix}$$

► Let's look at

$$H \cdot \mathbf{x} = \begin{bmatrix} \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \end{bmatrix} \cdot \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \\ x_7 \\ x_8 \end{bmatrix}$$



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We have

$$\begin{aligned}
 H \cdot \mathbf{x} &= \begin{bmatrix} \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \end{bmatrix} \cdot \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \\ x_7 \\ x_8 \end{bmatrix} \\
 &= \frac{1}{2} \begin{bmatrix} x_1 + x_2 \\ x_3 + x_4 \\ x_5 + x_6 \\ x_7 + x_8 \end{bmatrix} = \begin{bmatrix} x_1 & x_2 \\ x_3 & x_4 \\ x_5 & x_6 \\ x_7 & x_8 \end{bmatrix} \cdot \begin{bmatrix} 1/2 \\ 1/2 \end{bmatrix}
 \end{aligned}$$



- ▶ In a similar manner, we have

$$G \cdot \mathbf{x} = \begin{bmatrix} x_1 & x_2 \\ x_3 & x_4 \\ x_5 & x_6 \\ x_7 & x_8 \end{bmatrix} \cdot \begin{bmatrix} -1/2 \\ 1/2 \end{bmatrix}$$

- ▶ Thus to code $W_N \cdot \mathbf{x}$, we
 - ▶ Partition the input \mathbf{x} into

$$X = \begin{bmatrix} x_1 & x_2 \\ x_3 & x_4 \\ \vdots & \vdots \\ x_{N-1} & x_N \end{bmatrix}$$

- ▶ Compute $\mathbf{s} = X \cdot \begin{bmatrix} 1/2 \\ 1/2 \end{bmatrix}$, $\mathbf{d} = X \cdot \begin{bmatrix} -1/2 \\ 1/2 \end{bmatrix}$
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Here is a simple **Mathematica** module to do the job:

```
HWT1D[x_] := Module[{X, s, d, y},  
  X = Partition[x, 2, 2];  
  s = X.{1/2, 1/2};  
  d = X.{-1/2, 1/2};  
  y = Join[s, d];  
  Return[y];  
];
```



Here is a simple **Matlab** function:

```
function y=HWT1D(x)
    N=length(x)
    X=reshape(x,2,N/2);
    s=X. [.5; .5 ];
    d=X. [-.5; .5 ];
    y=[s; d];
```



The coding for the inverse is similar but with a different twist at the end. We need an algorithm for computing

$$W_8^{-1} \cdot \mathbf{y} = \left[\begin{array}{cccc|cccc} 1 & 0 & 0 & 0 & -1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & -1 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & -1 & 0 \\ 0 & 0 & 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & -1 \\ 0 & 0 & 0 & 1 & 0 & 0 & 0 & 1 \end{array} \right] \cdot \begin{bmatrix} y_1 \\ y_2 \\ y_3 \\ y_4 \\ y_5 \\ y_6 \\ y_7 \\ y_8 \end{bmatrix} = \begin{bmatrix} y_1 - y_5 \\ y_1 + y_5 \\ y_2 - y_6 \\ y_2 + y_6 \\ y_3 - y_7 \\ y_3 + y_7 \\ y_4 - y_8 \\ y_4 + y_8 \end{bmatrix}$$



- ▶ We could define the matrix

$$Y = \begin{bmatrix} y_1 & y_5 \\ y_2 & y_6 \\ y_3 & y_7 \\ y_4 & y_8 \end{bmatrix}$$

- ▶ and then compute

$$\mathbf{a} = Y \cdot \begin{bmatrix} 1 \\ -1 \end{bmatrix} = \begin{bmatrix} y_1 - y_5 \\ y_2 - y_6 \\ y_3 - y_7 \\ y_4 - y_8 \end{bmatrix}, \quad \mathbf{b} = Y \cdot \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} y_1 + y_5 \\ y_2 + y_6 \\ y_3 + y_7 \\ y_4 + y_8 \end{bmatrix}$$

- ▶ We return

$$(a_1, b_1, a_2, b_2, a_3, b_3, a_4, b_4)$$



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$$Y = \begin{bmatrix} y_1 & y_5 \\ y_2 & y_6 \\ y_3 & y_7 \\ y_4 & y_8 \end{bmatrix}$$

- ▶ and then compute

$$\mathbf{a} = Y \cdot \begin{bmatrix} 1 \\ -1 \end{bmatrix} = \begin{bmatrix} y_1 - y_5 \\ y_2 - y_6 \\ y_3 - y_7 \\ y_4 - y_8 \end{bmatrix}, \quad \mathbf{b} = Y \cdot \begin{bmatrix} 1 \\ 1 \end{bmatrix} = \begin{bmatrix} y_1 + y_5 \\ y_2 + y_6 \\ y_3 + y_7 \\ y_4 + y_8 \end{bmatrix}$$

- ▶ We return

$$(a_1, b_1, a_2, b_2, a_3, b_3, a_4, b_4)$$



Here is a simple **Mathematica** module to do the job:

```
IHWT1D[y_] := Module[{Y, a, b, x},  
  Y = Transpose[Partition[y, Length[y]/2]];  
  a = Y.{1, -1};  
  b = X.{1, 1};  
  x = Transpose[{a, b}];  
  Return[Flatten[x]];  
];
```



Let's have a look at the **Mathematica** notebook

```
HaarTransform1D.nb
```



An 8-bit digital image can be viewed as a matrix whose entries (known as *pixels*) range from 0 (black) to 255 (white).



129	128	121	51	127	224	201	179	159	140
148	116	130	75	184	191	182	185	186	180
175	169	166	195	195	192	168	173	166	158
157	171	169	182	199	205	191	191	180	172
73	89	96	100	122	143	166	190	188	180
93	107	103	81	70	77	106	139	165	181
106	105	112	132	144	147	189	183	158	184
102	100	106	124	140	157	179	175	168	175
91	105	112	93	86	85	100	104	110	106
97	97	112	102	113	111	105	94	103	104



Consider the 480×640 image (call it A)

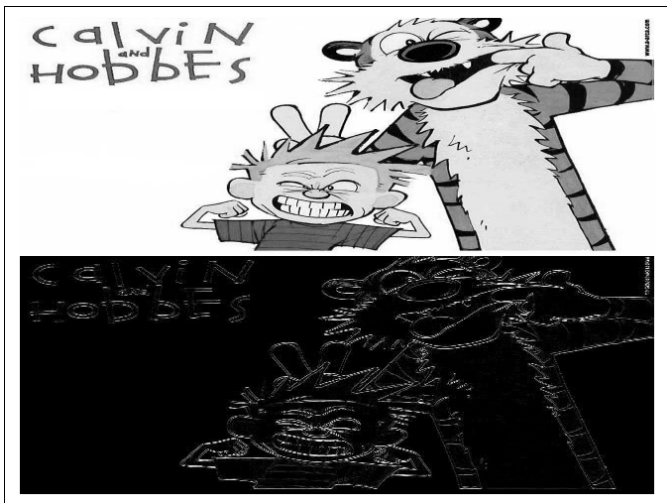


If $\mathbf{a}^1, \dots, \mathbf{a}^{640}$ are the columns of A , then computing $W_{640}A$ is the same as applying the HWT to each column of A :

$$W_{640}A = (W_{640} \cdot \mathbf{a}^1, \dots, W_{640} \cdot \mathbf{a}^{640})$$



Graphically, we have



- ▶ $C = W_{640}A$ processes the **columns** of A .
- ▶ How would we process the **rows** of C ?
- ▶ We compute $CW_{480}^T = W_{640}AW_{480}^T$.
- ▶ The **two-dimensional Haar transform** of $M \times N$ matrix A is

$$B = W_N A W_M^T$$



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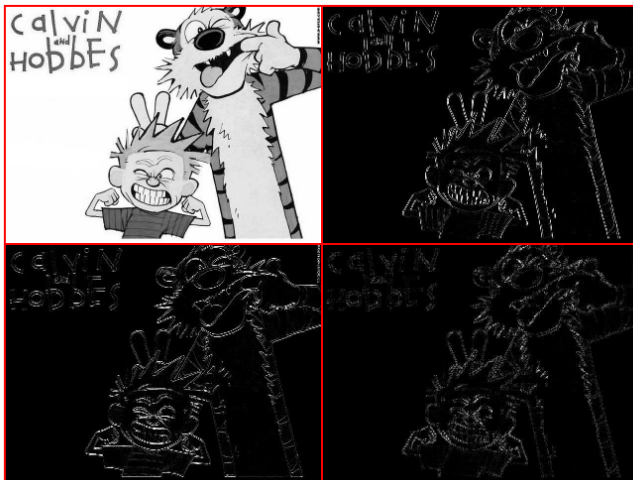


- ▶ $C = W_{640}A$ processes the **columns** of A .
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Graphically, we have



- ▶ Can we interpret what the transformation does to the image?
- ▶ Suppose A is the 4×4 matrix

$$A = \begin{bmatrix} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ a_{31} & a_{32} & a_{33} & a_{34} \\ a_{41} & a_{42} & a_{43} & a_{44} \end{bmatrix}$$

- ▶ Partitioning $W_4 = \begin{bmatrix} H \\ - \\ G \end{bmatrix}$, we have



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- ▶ Partitioning $W_4 = \begin{bmatrix} H \\ - \\ G \end{bmatrix}$, we have



$$\begin{aligned}
 W_4 A W_4^T &= \begin{bmatrix} H \\ - \\ G \end{bmatrix} A \begin{bmatrix} H^T & | & G^T \end{bmatrix} \\
 &= \begin{bmatrix} HA \\ - \\ GA \end{bmatrix} \begin{bmatrix} H^T & | & G^T \end{bmatrix} \\
 &= \left[\begin{array}{c|c} HAH^T & HAG^T \\ \hline GAH^T & GAG^T \end{array} \right]
 \end{aligned}$$

Let's look at each 2×2 block individually:



$$\blacktriangleright \mathit{HAH}^T = \frac{1}{4} \left[\begin{array}{cc|cc} \mathbf{a}_{11} + \mathbf{a}_{12} + \mathbf{a}_{21} + \mathbf{a}_{22} & & \mathbf{a}_{13} + \mathbf{a}_{14} + \mathbf{a}_{23} + \mathbf{a}_{24} & \\ & & & \\ \hline \mathbf{a}_{31} + \mathbf{a}_{32} + \mathbf{a}_{41} + \mathbf{a}_{42} & & \mathbf{a}_{33} + \mathbf{a}_{34} + \mathbf{a}_{43} + \mathbf{a}_{44} & \\ & & & \end{array} \right]$$

- ▶ Partition A in 2×2 blocks as

$$A = \left[\begin{array}{cc|cc} \mathbf{a}_{11} & \mathbf{a}_{12} & \mathbf{a}_{13} & \mathbf{a}_{14} \\ \mathbf{a}_{21} & \mathbf{a}_{22} & \mathbf{a}_{23} & \mathbf{a}_{24} \\ \hline \mathbf{a}_{31} & \mathbf{a}_{32} & \mathbf{a}_{33} & \mathbf{a}_{34} \\ \mathbf{a}_{41} & \mathbf{a}_{42} & \mathbf{a}_{43} & \mathbf{a}_{44} \end{array} \right] = \left[\begin{array}{c|c} \mathbf{A}_{11} & \mathbf{A}_{12} \\ \hline \mathbf{A}_{21} & \mathbf{A}_{22} \end{array} \right]$$

- ▶ Then the (i, j) element of HAH^T is simply the average of the elements in A_{ij} !
- ▶ So HAH^T is an approximation or **blur** of the original image. We will denote HAH^T as \mathcal{B} .



$$\blacktriangleright \mathit{HAH}^T = \frac{1}{4} \left[\begin{array}{cc|cc} a_{11} + a_{12} + a_{21} + a_{22} & a_{13} + a_{14} + a_{23} + a_{24} \\ a_{31} + a_{32} + a_{41} + a_{42} & a_{33} + a_{34} + a_{43} + a_{44} \end{array} \right]$$

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- ▶ The upper right hand corner is

$$HAG^T = \frac{1}{4} \begin{bmatrix} (a_{12} + a_{22}) - (a_{11} + a_{21}) & (a_{14} + a_{24}) - (a_{13} + a_{23}) \\ (a_{32} + a_{42}) - (a_{31} + a_{41}) & (a_{34} + a_{44}) - (a_{33} + a_{43}) \end{bmatrix}$$

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- ▶ The (i, j) element of HAG^T can be viewed as differences along columns of A_{ij} .
- ▶ We will denote HAG^T as \mathcal{V} (for vertical differences).



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$$HAG^T = \frac{1}{4} \begin{bmatrix} (a_{12} + a_{22}) - (a_{11} + a_{21}) & (a_{14} + a_{24}) - (a_{13} + a_{23}) \\ (a_{32} + a_{42}) - (a_{31} + a_{41}) & (a_{34} + a_{44}) - (a_{33} + a_{43}) \end{bmatrix}$$

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- ▶ The (i, j) element of HAG^T can be viewed as differences along columns of A_{ij} .
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- ▶ The lower left hand corner is

$$GAH^T = \frac{1}{4} \begin{bmatrix} (a_{21} + a_{22}) - (a_{12} + a_{11}) & (a_{23} + a_{24}) - (a_{13} + a_{14}) \\ (a_{31} + a_{32}) - (a_{42} + a_{41}) & (a_{43} + a_{44}) - (a_{33} + a_{34}) \end{bmatrix}$$

- ▶ Partition A in 2×2 blocks as

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- ▶ The (i, j) element of HAG^T can be viewed as differences along rows of A_{ij} .
- ▶ We will denote GAH^T as \mathcal{H} (for horizontal differences).



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- ▶ Partition A in 2×2 blocks as

$$A = \left[\begin{array}{cc|cc} a_{11} & a_{12} & a_{13} & a_{14} \\ a_{21} & a_{22} & a_{23} & a_{24} \\ \hline a_{31} & a_{32} & a_{33} & a_{34} \\ a_{41} & a_{42} & a_{43} & a_{44} \end{array} \right] = \left[\begin{array}{c|c} A_{11} & A_{12} \\ \hline A_{21} & A_{22} \end{array} \right]$$

- ▶ The (i, j) element of GAG^T can be viewed as differences along the diagonal of A_{ij} .
- ▶ We will denote GAG^T as \mathcal{D} (for diagonal differences).



- ▶ The lower right hand corner is

$$GAG^T = \frac{1}{4} \begin{bmatrix} (a_{11} + a_{22}) - (a_{12} + a_{21}) & (a_{13} + a_{24}) - (a_{23} + a_{14}) \\ (a_{31} + a_{42}) - (a_{32} + a_{41}) & (a_{33} + a_{44}) - (a_{43} + a_{34}) \end{bmatrix}$$

- ▶ Partition A in 2×2 blocks as

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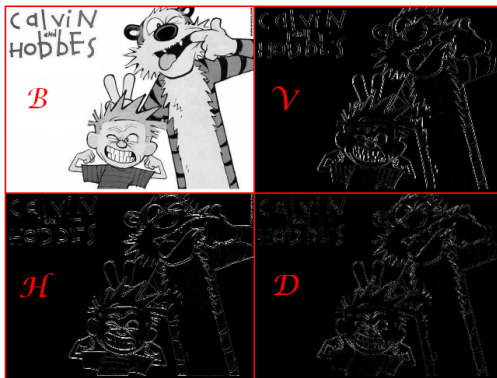
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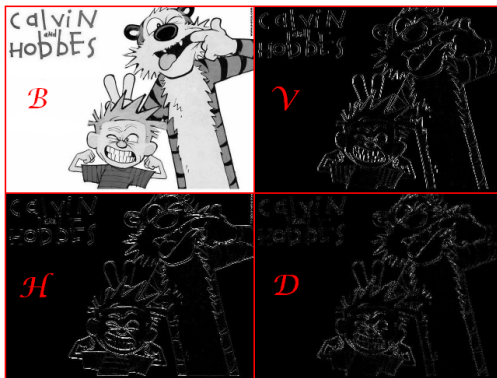
- ▶ So again the transform of our image is



- ▶ Once we have coded up the 2D HWT, I have the students write some Mathematica code to discover the parts of the transform.
- ▶ They can also think about how to build the inverse using the same code.



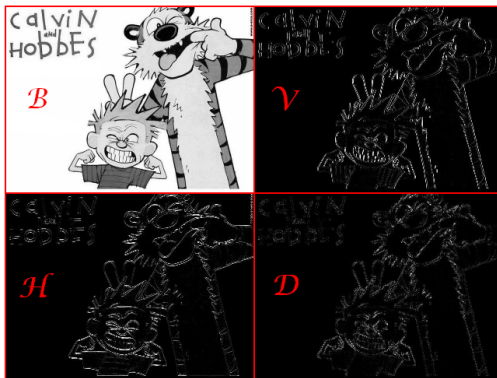
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- ▶ Writing code for the 2D Haar transform is easy and a good review of some basic properties from linear algebra.
- ▶ Given an $M \times N$ matrix A , we wish to compute

$$B = W_M A W_N^T$$

- ▶ We start by computing $C = W_M A$. This is easy - we simply apply HWT1D to each column of A .
- ▶ Mathematica is a “row-oriented” language so to perform $W_M A$, we must **transpose** A , apply HWT1D to A^T and then transpose back.
- ▶ Here is the Mathematica module:

```
LeftHaar[a_] := Transpose[Map[HWT1D, Transpose[a]]];
```



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Let's have a look at the Mathematica notebook

HaarTransform2D.nb



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- ▶ The process is quite simple.
 - ▶ Compute i iterations of the HWT on A .
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HaarEdgeDetection.nb



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 - ▶ Perform i iterations of the HWT on A .
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- ▶ **Cumulative Energy** is a simple vector-valued function that gives information about the concentration of energy in a vector.
- ▶ Let $\mathbf{v} = (v_1, \dots, v_N) \in \mathbb{R}^N$.
 - ▶ Take the absolute value of each component of \mathbf{v} and sort from largest to smallest. Call this new vector \mathbf{y} .
 - ▶ Note $\|\mathbf{v}\| = \|\mathbf{y}\|$.
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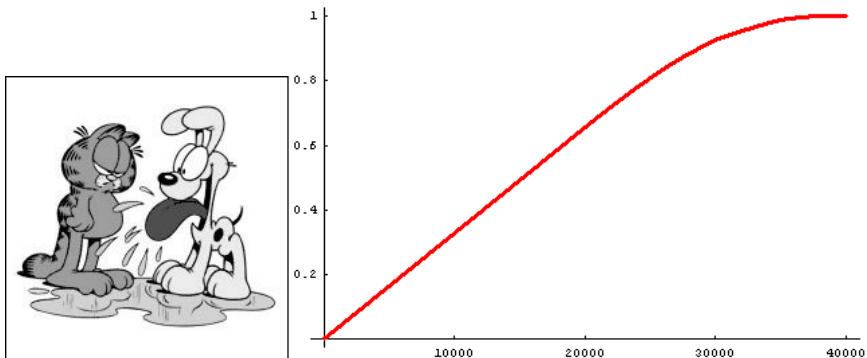
$$CE_k = \sum_{j=1}^k \frac{y_j^2}{\|\mathbf{y}\|^2}$$



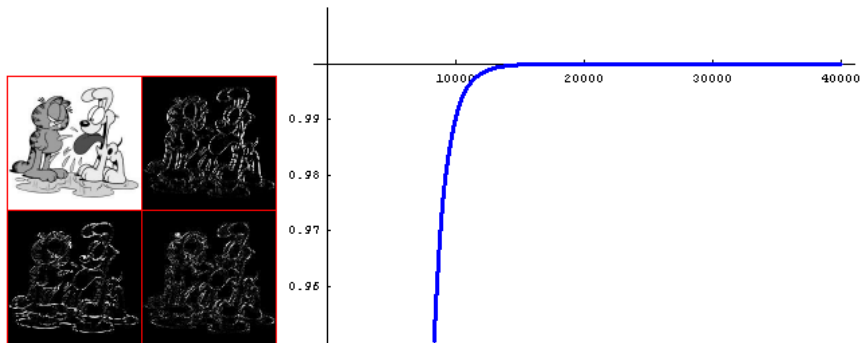
$$\begin{aligned}CE_1 &= \frac{y_1^2}{y_1^2 + y_2^2 + \cdots + y_N^2} \\CE_2 &= \frac{y_1^2 + y_2^2}{y_1^2 + y_2^2 + \cdots + y_N^2} \\&\vdots \\CE_N &= \frac{y_1^2 + y_2^2 + \cdots + y_N^2}{y_1^2 + y_2^2 + \cdots + y_N^2} = 1\end{aligned}$$



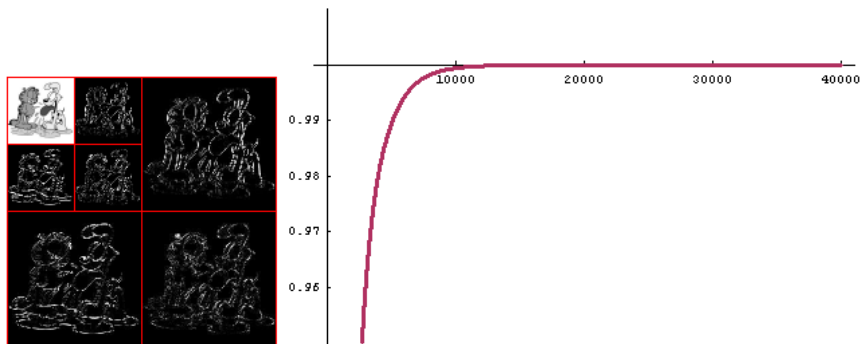
Here is the cumulative energy plot for the Garfield image:



Here is the cumulative energy plot for one iteration of the transformed image:



Here is the cumulative energy plot for two iterations of the transformed image:



We will use Cumulative Energy to

- ▶ Determining an amount of energy we wish to retain.
- ▶ Identifying those elements in the vector that produce that energy.
- ▶ Converting the remaining elements to 0.



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- ▶ In 1952, **David Huffman** made a simple observation:
- ▶ *Rather than use the same number of bits to represent each character, why not use a short bit stream for characters that appear often in an image and a longer bit stream for characters that appear infrequently in the image?*
- ▶ He then developed an algorithm to do just that. We refer to his simple algorithm as **Huffman encoding**. We will illustrate the algorithm via an example.



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- ▶ Suppose you want to perform Huffman encoding on the word **seesaws**.
- ▶ First observe that *s* appears three times (24 bits), *e* appears twice (16 bits), and *a* and *w* each appear once (16 bits) so the total number of bits needed to represent *seesaws* is 56.



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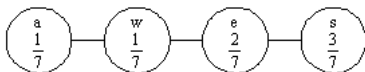
Char.	ASCII	Binary	Frequency
s	115	01110011 ₂	3
e	101	01100101 ₂	2
a	97	01100001 ₂	1
w	119	01110111 ₂	1

So in terms of bits, the word **seesaws** is

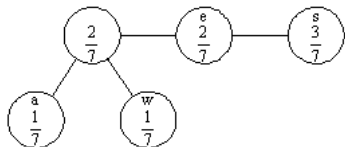
01110011 01100101 01100101 01110011 01100001 01110111 01110011



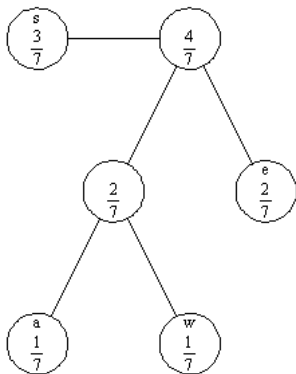
The first step in Huffman coding is as follows: Assign probabilities to each character and then sort from smallest to largest. We will put the probabilities in circles called **nodes** and connect them with lines (**branches**).



Now simply add the two smallest probabilities to create a new node with probability $2/7$. Branch the two small nodes off this one and resort the three remaining nodes:

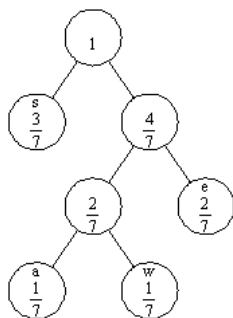


Again we add the smallest two probabilities on the top row ($2/7 + 2/7 = 4/7$), create a new node with everything below these nodes as branches and sort again:



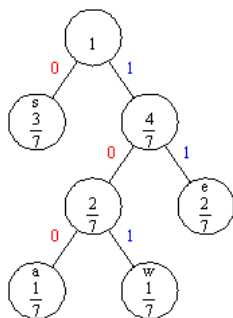
Since only two nodes remain on top, we simply add the probabilities of these nodes together to get 1 and obtain our finished tree:





Now assign to each left branch the value 0 and to each right branch the value 1:





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- ▶ Here are the new bit streams for the four characters:



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- ▶ Here are the new bit streams for the four characters:



Char.	Binary
s	0_2
e	11_2
a	100_2
w	101_2



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- ▶ The total number of bits we need to represent the word *seesaws* is 13 bits! Recall without Huffman coding, we needed 56 bits so we have reduced the number of bits needed by a factor of 4!
- ▶ Here is the word *seesaws* using the Huffman codes for each character:

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We are ready to perform naive image compression. Let's have a look at the notebook

```
HaarImageCompression.nb
```

